

# CS 13: Mathematical Foundations of Computing

## Lecture Exercises

Name:

E-mail:

0.

### Functions on Trees

$$\text{flip}(\text{Nil}) = \text{Nil}$$

$$\text{flip}(\text{Tree}(x, L, R)) = \text{Tree}(x, \text{flip}(R), \text{flip}(L))$$

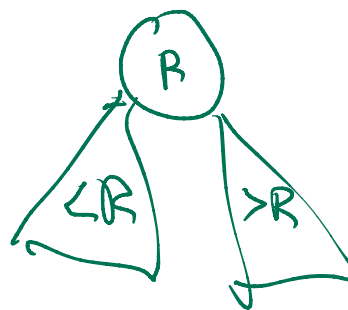
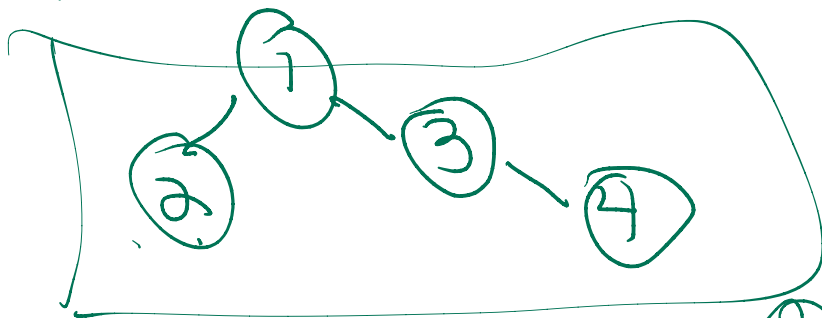
### Claim

For all  $T \in \text{Tree}$ ,

$$\text{flip}(\text{flip}(T)) = T$$

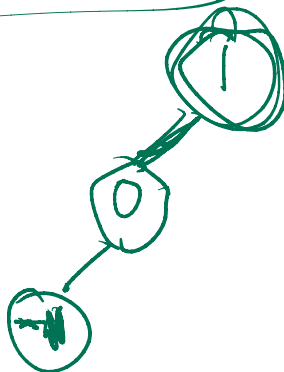
$\text{Tree} = \underline{\text{Nil}} \mid \text{Tree}(x, L, R)$  for  $x \in \mathbb{Z}, L, R \in \text{Tree}$

$\text{Tree}(1, \text{Tree}(2, \text{Nil}, \text{Nil}), \text{Tree}(3, \text{Nil}, \text{Tree}(4, \text{Nil}, \text{Nil})))$



case Nil:

case  $T(x, L, R)$



## 1.

### Functions on Trees

$\text{full}(0) = \text{Nil}$

$\text{full}(n) = \text{Tree}(n, \text{full}(n-1), \text{full}(n-1))$

$\text{num\_leaves}(\text{Nil}) = 1$

$\text{num\_leaves}(\text{Tree}(x, L, R)) = \text{num\_leaves}(L) + \text{num\_leaves}(R)$

### Claim

For all  $n \in \mathbb{N}$ ,

$$\text{num\_leaves}(\text{full}(n)) = 2^n$$

2.

Functions on Trees

$$\begin{aligned} \text{size}(\text{Nil}) &= 0 \\ \text{size}(\text{Tree}(x, L, R)) &= 1 + \text{size}(L) + \text{size}(R) \\ \text{height}(\text{Nil}) &= 0 \\ \text{height}(\text{Tree}(x, L, R)) &= 1 + \max(\text{height}(L), \text{height}(R)) \end{aligned}$$

Claim

For all  $T \in \text{Tree}$ ,

$$\text{size}(T) \leq 2^{\text{height}(T)} - 1$$

Handwritten proof:

$$\begin{aligned} \text{size}(T(x, L, R)) &= 1 + \text{size}(L) + \text{size}(R) \\ &\leq 1 + 2^{\text{height}(L)} - 1 + 2^{\text{height}(R)} - 1 \quad \text{by IH} \\ &\leq 1 + 2^{\max(\text{height}(L), \text{height}(R))} - 2 \quad \text{This step is non-obvious} \\ &= 2^{\max(\text{height}(L), \text{height}(R))} - 1 \\ &= 2^{\text{height}(T(x, L, R))} - 1 \end{aligned}$$

Additional notes:  $2^{\text{height}(L)} \leq 2^{\max(\text{height}(L), \text{height}(R))}$